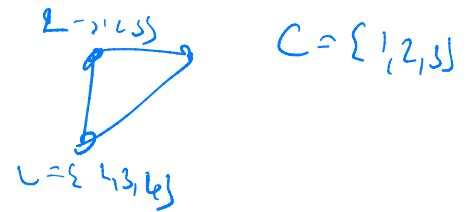


Lecture 8: List Coloring

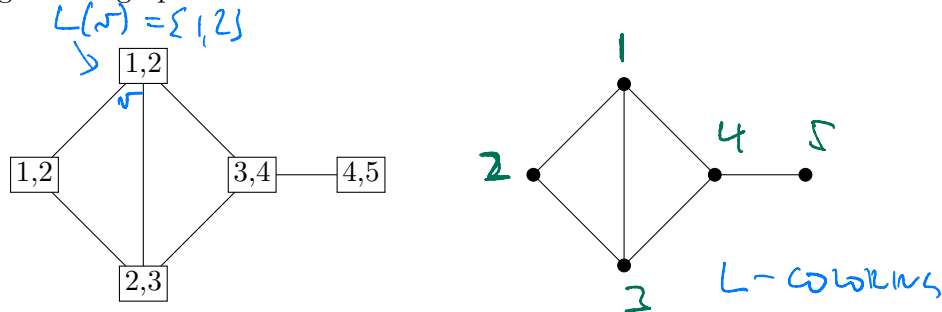
Introduced by Vizing (1976) and independently by Erdős, Rubin, and Taylor (1979).

A *list assignment* of G is a function L that assigns to each vertex $v \in V(G)$ a list $L(v)$ of colors. The elements of the list $L(v)$ are called *admissible colors* for the vertex v . An L -*coloring* is a mapping $\varphi : V(G) \rightarrow \bigcup_v L(v)$ such that

- $\varphi(v) \in L(v)$ for every $v \in V(G)$, and
- $\varphi(u) \neq \varphi(v)$ whenever u and v are adjacent vertices of G .



1: Find a list coloring for the graph below.



For a graph G we say

- L -colorable if G admits an L -coloring
- k -choosable if, for every list assignment L with $|L(v)| \geq k$ for all $v \in V(G)$, G is L -colorable.
- *choosability* of a graph G , denoted by $ch(G)$ or $\chi_\ell(G)$, is the smallest k such that G is k -choosable.

2: What is choosability of the following graphs? (Some depicted twice to allow experiments.)

$\chi_\ell(K_5) \leq \chi_\ell(K_4)$

K_5 is NOT 2-choosable

$\chi_\ell(K_5) = 3$

$\chi_\ell(K_4) = 2 \leq \chi_\ell(K_5)$

$\chi_\ell(K_4) = 3$

$\chi_\ell(K_4) = 2$

Erdős, Rubin and Taylor showed that there are bipartite graphs with arbitrary large list chromatic number.

Theorem 1 (Erdős, Rubin and Taylor). For $m \geq \binom{2k-1}{k}$, the bipartite graph $K_{m,m}$ is not k -choosable.

3: Prove the theorem. Hint: Use the last graph from previous exercise.

$\{1, 2, \dots, k\}$

$(1, 2, \dots, 2k-1)$

$L(v) \in \binom{\{1, 2, \dots, k\}}{2k-1}$

AND USE ALL SUBSETS IN EACH A & B

IN L-COLORING AT LEAST k COLORED ON B (DISTINCT)

① 2 3 4 5 6 7 8

1 2, 3, 4 | $2k$

1, ..., k-1 | $\ominus k$ | $2k-1$

$\{1, 2, \dots, k\}$
 $\{2k-1, \dots, 2k\}$
 $\{2k, \dots, 2k+1\}$
 \dots

ALL SUBSETS OF $\{1, \dots, k\}$
 OF SIZE $2k-1$
 OF SIZE $2k$
 OF SIZE $2k+1$
 \dots

Let G be a graph from which we start removing vertices of degree one consecutively one by one. By this procedure, we end up either by a vertex of minimum degree ≥ 2 or by a single vertex. We denote the resulting graph by $\text{core}(G)$. Recall that the Theta graph $\Theta_{a,b,c}$ is comprised from two vertices that are connected by three paths of length a , b , and c that are pair-wise disjoint except at the end-vertices.

Theorem 2 (Rubin). A graph G is 2-choosable if and only if

$$\text{core}(G) \in \{K_1, C_{2m+2}, \Theta_{2,2,2m} : m \geq 1\}.$$

4: Show that

$$\text{core}(G) \in \{K_1, C_{2m+2}, \Theta_{2,2,2m} : m \geq 1\}$$

are 2-choosable.

1 List version of Brooks theorem and beyond

5: Show that

$$\chi_\ell(G) \leq \Delta(G) + 1.$$

Moreover, k -degenerated graph is $(k+1)$ -choosable.

$\Delta(G) + 1$

Colors

v_1, v_2, v_3, v_4

GREEDY COLORING WORKS

k

6: State Brook's theorem.

$$\chi(G) = \Delta(G) + 1 \text{ IFF } G \text{ is } K_n \text{ or } C_{2n+1} \text{ for } n \geq 1$$

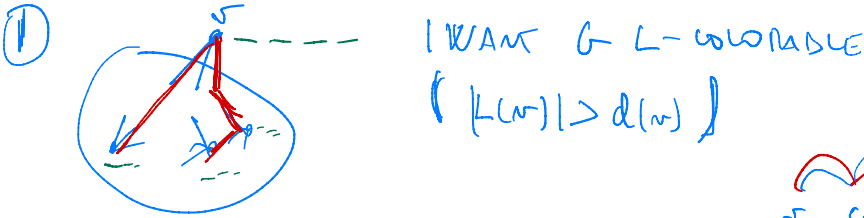


Question: When is graph L -colorable under assumption that $|L(v)| \geq d(v)$ for every vertex v ? We call such an assignment a *degree list assignment*.

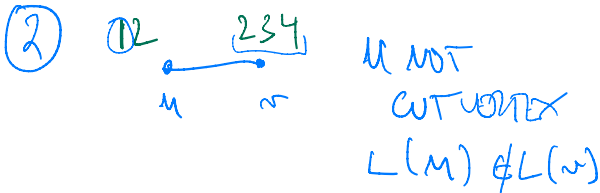
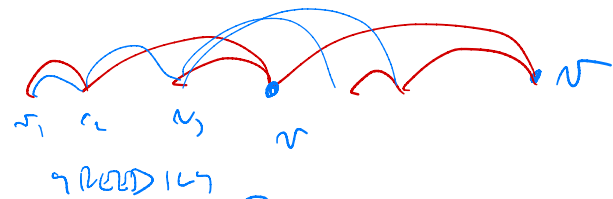
Lemma 3. Let (G, L) be a pair of a connected graph and L a degree list assignment such that G is not L -colorable. Then following hold:

- (1). $|L(v)| = d(v)$ for every vertex v of G ;
- (2). If u and v are two adjacent vertices of G and u is not a cut-vertex then $L(u) \subseteq L(v)$;
- (3). If G is 2-connected then it is an odd cycle or a complete graph and L assigns the same $\Delta(G)$ colors to all vertices.

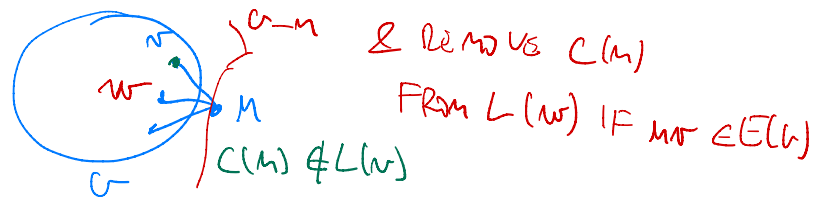
7: Prove the lemma.



SPANNING TREE



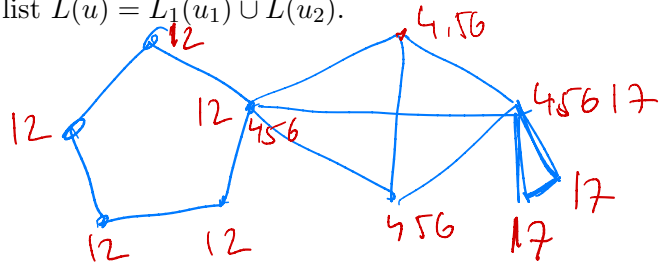
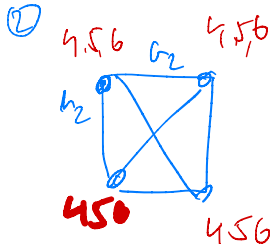
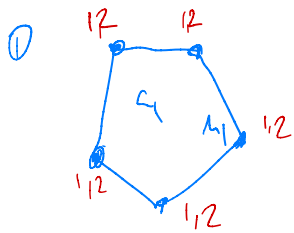
THEN G IS L -COLORABLE
 $\Rightarrow d(u) \leq d(v)$



③ G IS d -REGULAR
 & ALL LISTS ARE THE SAME
 \Rightarrow BROOKS $\Rightarrow C_{2k+1}$ OR K_n
 WHEN NOT Δ -COLORABLE
 BUT $|L_{u-v}(u)| = |L(u)| > d_{u-v}(u)$

$u, v \in G - u \dots d_{u-v}(u) = d(u) - 1$
 $|L_{u-v}(u)| = |L(u)| > d_{u-v}(u)$
 $|L_{u-v}(u)| \geq |L(u)| - 1$

1. G is an odd cycle and L assigns a same two colors to every vertex of G ; or
2. G is a complete graph and L assigns a same colors to every vertex of G ; or
3. there exists pairs (G_1, L_1) and (G_2, L_2) from \mathcal{G}_L such that G is obtained from identifying a vertex u_1 from G_1 with a vertex u_2 from G_2 into a vertex u of G , where $L_1(u_1) \cap L_2(u_2) = \emptyset$. And, L coincides with L_1 and L_2 except at the identified vertex u it has list $L(u) = L_1(u_1) \cup L_2(u_2)$.

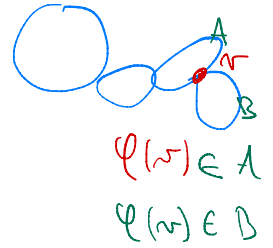


Theorem 4. A connected graph G is not L -colorable for a degree list assignment L if and only if $(G, L) \in \mathcal{G}_L$.

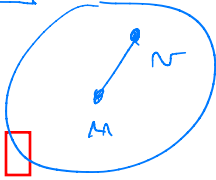
8: Hint: Prove the theorem.

IF $(G, L) \notin \mathcal{G}_L$ THEN G NOT L -COLORABLE

IS \neq OF 2-CONNECTED COMPONENTS OPPOSE $\exists \varphi$ L -EQUIVARIANT
 NOT L -COLORABLE \Rightarrow GALWAY TREE



LEMMA 3

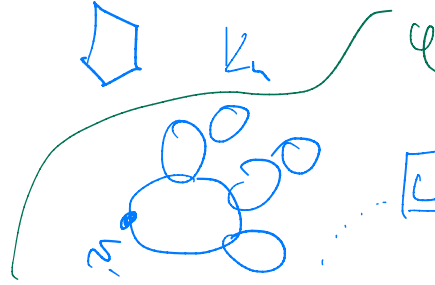


G 2-CONNECTED

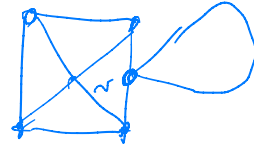
LEMMA 3 \rightarrow



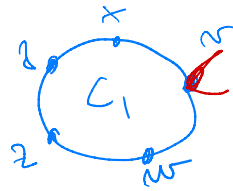
$C-N$



$L(m)$



$C \in L(m)$
 $\varphi(m) = C$



$L(x) \leq L(y)$
 $L(y) = L(x) \Leftrightarrow d(x) = d(y)$
 $L(x) \leq L(y)$
 $L(y) \leq L(x)$ } $L \geq 2$

$\forall x, y \in C_1 \setminus \{m\}$
 $L(x) = L(y) \in L(m)$
 $d(x) = d(y)$
 $L(y) = \widehat{2}$

This asserts our first proposition.

Corollary 5. For any graph G that is not an odd cycle or a complete graph, holds

$$\chi_L(G) \leq \Delta(G).$$

9: Prove the corollary.

LET G HAVE LIST ASSIGNMENT L $|L(w)| = \Delta(w) \forall w$

& G NOT L -COLORABLE

IF $\forall w |L(w)| \geq d(w)$ L 3-APPLIES

\rightarrow IF 2-CONNECTED \Rightarrow ODD CYCLE OR COMPLETE GRAPH



L 3.1 G IS Δ -REGULAR & GALWAY TREE IS MUST BE 2-CONNECTED

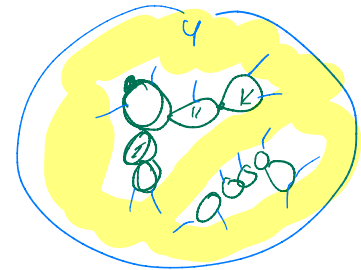
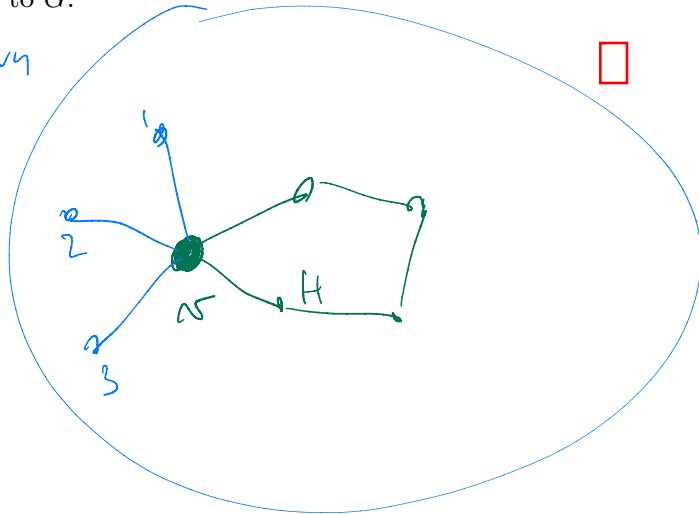
Theorem 6 (Gallai). In a k -critical graph with $k \geq 4$, low vertices induce a forest (possibly empty) whose blocks are odd cycles and complete graphs.

Recall that low vertices are vertices of degree $k - 1$.

k -CRITICAL GRAPH

10: Prove the theorem. Hint: Consider G without low vertices L and try to extend $(k - 1)$ coloring of $G - L$ to G .

TRYING TO $k-1$ COLOR



k -CRITICAL GRAPH G
 $\delta(G) \geq \boxed{k-1}$

$d(v) = k - 1$

" $L(v) = k - 1$ "

$L(v) = \{1, \dots, k-1\} \setminus \{\varphi(x), x \in N(v) \cap V(G-L)\}$

$\Rightarrow L(v) \geq d_H(v)$ BUT \downarrow CANNOT COLOR $H \Rightarrow H$ IS ACYCLIC TREE

$d(v) = k + 2$
 " $L(v) = k - 1$ "

\Rightarrow

$d(v) = k + 1$
 $L(v) = k - 2$

$d(v) \not\leq |L(v)|$

1.1 Planar graphs

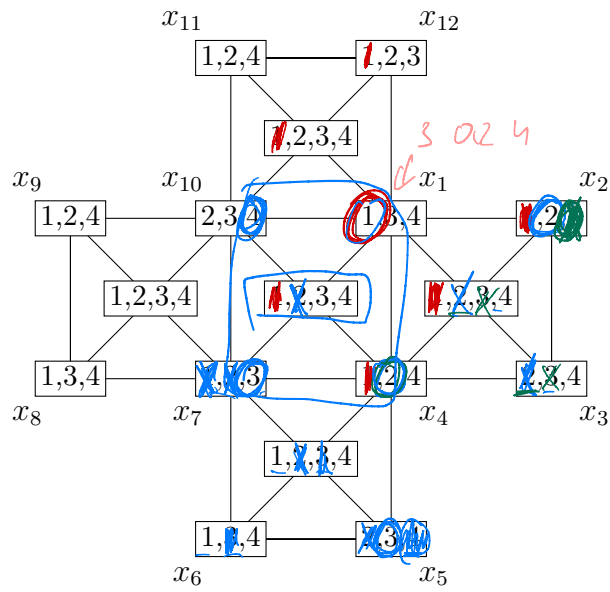
Theorem 7 (Thomassen). *Every planar graph is 5-choosable.*

The next lemma implies the above theorem. Recall that a near-triangulation is a plane graph whose all inner faces are 3-cycles.

Lemma 8. *Let G be a 2-connected near-triangulation and let $C = x_1x_2 \cdots x_nx_1$ be the outerface. Let L be a list-assignment of G such that $|L(x)| \geq 3$, for $x \in V(C)$, and otherwise $|L(x)| \geq 5$. Suppose that c is an L -coloring of x_1 and x_n . Then, c can be extended to an L -coloring of G .*

11: Prove the lemma by induction.

Voigt construct a non-4-choosable planar graph on 238 vertices. Later Mirzakhani (the famous one) such a graph on 63 vertices. A gadget of her construction is depicted below.



12: Show that the graph above is not list-colorable and the graph on the next page is also not list-colorable.

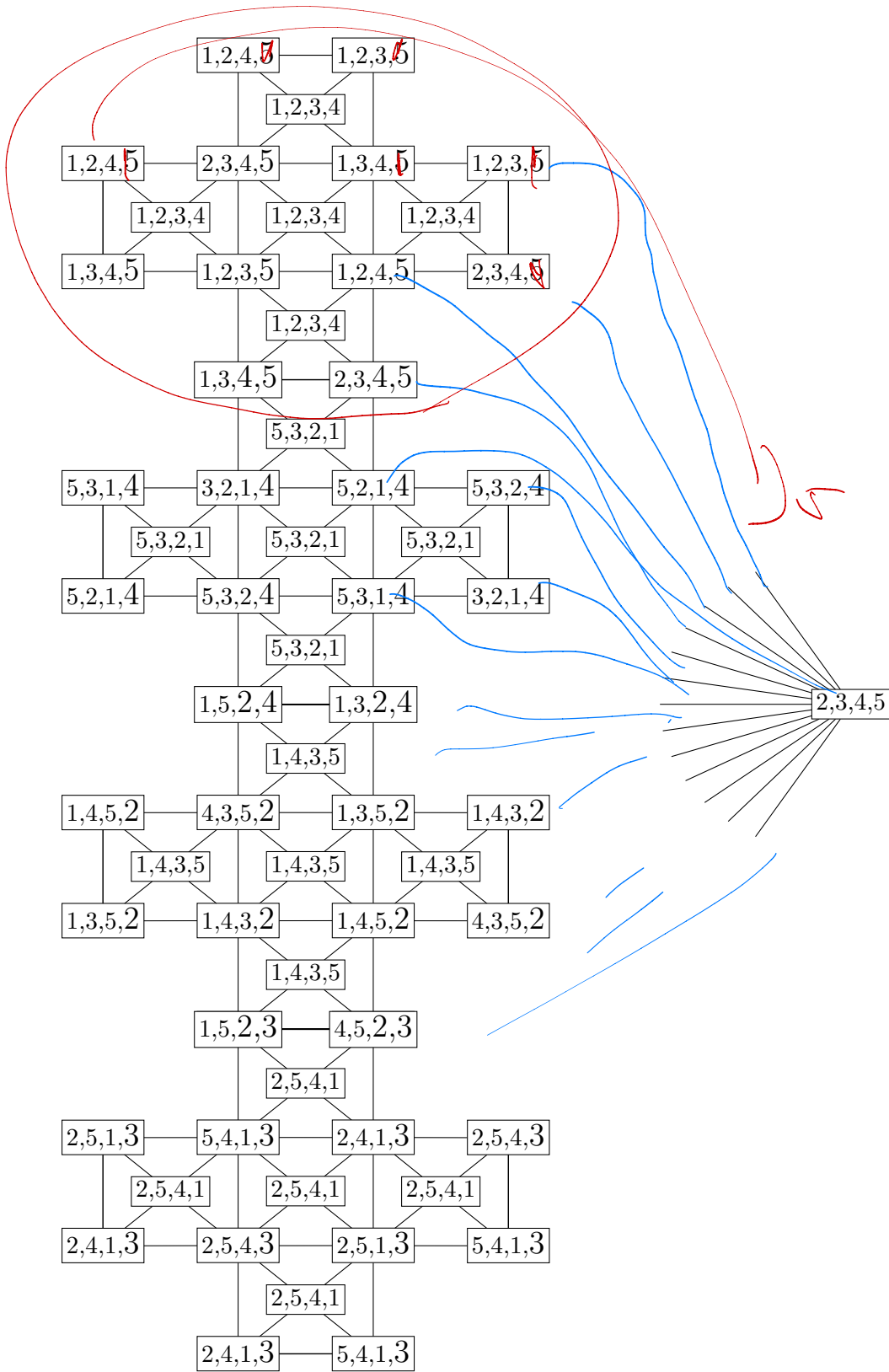


Figure 1: Mirzakhani construction.